

Introduction to OpenACC 2021 HPC Workshop: Parallel Programming

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Research Computing

July 13 - 15, 2021



- CPU : consists of a few cores optimized for sequential serial processing
- GPU : has a massively parallel architecture consisting of thousands of smaller, more efficient cores designed for handling multiple tasks simultaneously GPU enabled applications





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GPU enabled applications





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GPU Design



GPU Design

Model GPU design

- Large number of cores working in SIMD mode.
- Slow global memory access, high bandwidth.
- CPU communication over PCI bus.
- Warp scheduling and fast switching queue model.



Hetergenous Programming





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GPU Programming Languages



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Accelerator Fundamentals



• We must expose enough parallelism to saturate the device

- Accelerator threads are slower than CPU threads
- Accelerators have orders of magnitude more threads
- Fine grained parallelism is good
- Coarse grained parallelism is bad
 - Lots of legacy apps have only exposed coarse grain parallelism
 - i.e. MPI and possibly OpenMP

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What is OpenACC?





- Open Standard
- Easy, Compiler-Driven Approach
- portable across host CPUs and accelerators

What is OpenACC?



History

- OpenACC was developed by The Portland Group (PGI), Cray, CAPS and NVIDIA.
- PGI, Cray, and CAPs have spent over 2 years developing and shipping commercial compilers that use directives to enable GPU acceleration as core technology.
- The small differences between their approaches allowed the formation of a group to standardize a single directives approach for accelerators and CPUs.
- Full OpenACC 3.1 Specification available online: https://www.openacc.org/specification

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The Standard for GPU Directives

- Simple: Directive are the easy path to accelerate compute intensive applications
 - **Open:** OpenACC is an open GPU directives standard, making GPU programming straightforwards and portable across parallel and multi-core processors
- **Powerful:** GPU directives allow complete access to the massive parallel power of a GPU

What is OpenACC?



High Level

- Compiler directives to specify parallel regions in C & Fortran
 - Offload parallel regions
 - Portable across OSes, host CPUs, accelerators, and compilers
- Create high-level heterogenous programs
 - Without explicit accelerator intialization
 - Without explicit data or program transfers between host and accelerator

High Level · · · with low-level access

- Programming model allows programmers to start simple
- Compiler gives additional guidance
 - Loop mappings, data location and other performance details
- Compatible with other GPU languages and libraries
 - Interoperate between CUDA C/Fortran and GPU libraries
 - e.g. CUFFT, CUBLAS, CUSPARSE, etc

Why OpenACC



- Directives are easy and powerful.
- Avoid restructuring of existing code for production applications.
- Focus on expressing parallelism.

OpenACC is not GPU Programming

OpenACC is Expressing Parallelism in your code

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OpenACC Execution Model



- Application code runs on the CPU (sequential, shared or distributed memory)
- OpenACC directives indicate that the following block of compute intensive code needs to be offloaded to the GPU or accelerator.

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С/С++

#pragma acc directive [clause [,] clause] ...] ... often followed by a structured code block

Fortran

!\$acc directive [clause [,] clause] ...]

...often paired with a matching end directive surrounding a structured code block:

!\$acc end directive



C: #pragma acc kernels [clause]

Fortran !\$acc kernels [clause]

- The kernels directive expresses that a region may contain parallelism and the compiler determines what can be safely parallelized.
- The compiler breaks code in the kernel region into a sequence of kernels for execution on the accelerator device.
- What is a kernel? A function that runs in parallel on the GPU.
- When a program encounters a kernels contruct, it will launch a sequence of kernels in order on the device.

Fortran

```
!Sacc kernels
do i = 1, n
x(i) = 1.0
y(i) = 2.0
end do
do i = 1, n
y(i) = y(i) + a * x(i)
end do
```

!\$acc end kernels

С/С++

```
#progma acc kernels
{
  for (i = 0; i < n; i++){
    x[i] = 1.0;
    y[i] = 2.0;
  }
  for (i = 0; i < n; i++){
    y[i] = 0*x[i] + y[i];
  }
}</pre>
```



• The parallel directive identifies a block of code as having parallelism.

- Compiler generates a parallel kernel for that loop.
- C: #pragma acc parallel [clauses]

Fortran: !\$acc parallel [clauses]

Fortran

```
 \begin{array}{l} \text{!Sacc parallel} \\ \text{do } i = 1, n \\ \times(i) = 1.0 \\ y(i) = 2.0 \\ \text{end } \text{do} \\ \text{do } i = 1, n \\ y(i) = y(i) + a * x(i) \\ \text{end } \text{do} \\ \text{!Sacc end parallel} \end{array}
```

```
C/C++
```

```
#pragma acc parallel
{
  for (i = 0; i < n; i++){
    x[i] = 1.0;
    y[i] = 2.0;
  }
  for (i = 0; i < n; i++){
    y[i] = a*x[i] + y[i];
  }
}</pre>
```

OpenACC Loop Directive



- Loops are the most likely targets for Parallelizing.
- The Loop directive is used within a parallel or kernels directive identifying a loop that can be executed on the accelerator device.
- C: #pragma acc loop [clauses]
- Fortran: !\$acc loop [clauses]
 - The loop directive can be combined with the enclosing parallel or kernels
 - C: #pragma acc kernels loop [clauses]
- Fortran: !\$acc parallel loop [clauses]
 - The loop directive clauses can be used to optimize the code. This however requires knowledge of the accelerator device.
- Clauses: gang, worker, vector, num_gangs, num_workers

Fortran	C/C++	
<pre>!Sacc loop do i = 1, n y(i) = y(i) + a * x(i) end do !Sacc end loop</pre>	<pre>#progma acc loop for (i = 0; i < n; i++){ y[i] = a**[i] + y[i]; }</pre>	
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PARALLEL

- Requires analysis by programmer to ensure safe parallelism.
- Straightforward path from OpenMP

KERNELS

- Compiler performs parallel analysis and parallelizes what it believes is safe.
- Can cover larger area of code with single directive
- Gives compiler additional leeway

Both approaches are equally valid and can perform equally well.

Exercise



- Parallelize the saxpy code by adding OpenACC directives parallel or kernels
- Compile the code using the following flags to the NVIDIA HPC SDK compiler

-acc -gpu=cc75

• Run the code and compare timing with serial and openmp code

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Serial Code

```
program saxpy
 implicit none
 integer :: i,n
 real,dimension(:),allocatable :: x, y
 real :: a,start_time, end_time
 n = 200000000
 allocate(x(n),y(n))
  x = 1.000
 v = 2.000
  a = 2.0d0
 call cpu_time(start_time)
 do i = 1, n
     y(i) = y(i) + a * x(i)
  end do
 call cpu_time(end_time)
 deallocate(x,y)
 print '(a,f8.6)', 'SAXPY Time: ', end_time - start_time
end program saxpy
```

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SAXPY: OpenMP



OpenMP Code

program saxpy

```
implicit none
 integer :: i,n,omp_get_max_threads
 real,dimension(:),allocatable :: x, y
 real :: a,start_time, end_time
 n = 200000000
 allocate(x(n), y(n))
 !$omp parallel do default(shared) private(i)
 do i = 1, n
     x(i) = 1.0
    y(i) = 1.0
 end do
 !$omp end parallel do
  a = 2.0
 call cpu_time(start_time)
 !$omp parallel do default(shared) private(i)
 do i = 1, n
     y(i) = y(i) + a + x(i)
  end do
 !Somp end parallel do
 call cpu_time(end_time)
 deallocate(x,y)
 print '(a,i3,2x,a,f8.6)', 'Threads: ', omp_get_max_threads(), &
         'SAXPY Time: '. end time - start time
end program saxpy
```

```
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```

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SAXPY: OpenACC



OpenACC Code

program saxpy

```
implicit none
 integer :: i,n
 real,dimension(:),allocatable :: x, y
 real :: a,start_time, end_time
  n = 200000000
 allocate(x(n),y(n))
 !$acc parallel loop
 do i = 1. n
     x(i) = 1.0
    y(i) = 1.0
  end do
 !$acc end parallel loop
 a = 2.0
 call cpu_time(start_time)
 !$acc parallel loop
 do i = 1. n
     y(i) = y(i) + a * x(i)
  end do
 !Sacc end parallel loop
 call cpu_time(end_time)
 deallocate(x, y)
 print '(a,f8.6)', 'SAXPY Time: ', end_time - start_time
end program saxpy
```

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SAXPY: CUDA Fortran



CUDA Fortran Code

```
module mymodule
contains
 attributes(global) subroutine saxpy(n, a, x, y)
    real :: x(:), y(:), a
   integer :: n, i
   attributes(value) :: a, n
    i = threadIdx%x+(blockIdx%x-1)*blockDim%x
   if (i \le n) y(i) = a x(i) + y(i)
 end subroutine saxpy
end module mymodule
program main
 use cudafor; use mymodule
 integer, parameter :: n = 20000000
 real, device :: x_d(n), y_d(n)
 real. device :: a d
 real :: start_time, end_time
  x d = 1.0
  v d = 2.0
 a d = 2.0
 call cpu time(start time)
 call saxpy<<<4096, 256>>>(n, a, x_d, y_d)
 call cpu time(end time)
 print '(a,f15.6,a)', 'SAXPY Time: ', end_time - start_time, 'in secs'
end proaram main
```

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SAXPY: Compile & Run



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Compile

[alp514.sol](1018): nvc -acc -gpu=cc75 -o saxpyc_acc saxpy_acc.c

- Specify the gpu architecture: -gpu=cc75
- Get more information about the compilation: -Minfo=accel

```
[alp514.hawk-b625](1005): nvfortran -acc -gpu=cc75 -Minfo=accel -o saxpyf_acc saxpy_acc.f90
saxpy:
12, Generating Tesla code
13, ISacc loop vector(128) ! threadidx%x
12, Generating implicit copyout(x(1:200000000)) [if not already present]
13, Loop is parallelizable
15, Generating Tesla code
16, ISacc loop vector(128) ! threadidx%x
15, Generating implicit copyout(y(1:200000000)) [if not already present]
16, Loop is parallelizable
20, Generating Tesla code
21, ISacc loop gang, vector(128) ! blockidx%x threadidx%x
20, Generating implicit copy(y(1:200000000)) [if not already present]
Generating implicit copy(x(1:200000000)) [if not already present]
30, Generating implicit copy(x(1:200000000)) [if not already present]
31, ISacc loop gang vector(128) ! blockidx%x threadidx%x
32, Generating implicit copy(x(1:200000000)) [if not already present]
32, ISacc loop gang vector(128) ! blockidx%x threadidx%x
34, Generating implicit copy(x(1:200000000)) [if not already present]
34, ISacc loop gang vector(128) ! blockidx%x threadidx%x
34, Generating implicit copy(x(1:200000000)) [if not already present]
34, ISacc loop gang vector(128) ! blockidx%x threadidx%x
34, Generating implicit copy(x(1:200000000)) [if not already present]
34, ISacc loop gang vector(128) ! blockidx%x threadidx%x
34, Generating implicit copy(x(1:200000000)) [if not already present]
34, ISacc loop gang vector(128) ! ISacc loop gang vector(128) ! blockidx%x threadidx%x
34, Generating implicit copy(x(1:200000000)) [if not already present]
34, ISacc loop gang vector(128) ! ISacc loop g
```



Fortran Timings

Algorithm	Device	Time (s)	Speedup
Serial	Xeon Gold 5220R	0.504534	
OpenMP (12 threads)	Xeon Gold 5220R	0.050300	10.03
OpenMP (24 threads)	Xeon Gold 5220R	0.026623	18.95
OpenMP (48 threads)	Xeon Gold 5220R	0.025263	19.97
OpenACC	Tesla T4	0.517426	0.98
CUDA	Tesla T4	0.007846	64.30

C Timings

Algorithm	Device	Time (s)	Speedup
Serial	Xeon Gold 5220R	0.512128	
OpenMP (12 threads)	Xeon Gold 5220R	0.056454	9.07
OpenMP (24 threads)	Xeon Gold 5220R	0.048442	10.57
OpenMP (48 threads)	Xeon Gold 5220R	0.026348	19.44
OpenACC	Tesla T4	3.434997	0.15

What's going with OpenACC code? Why even bother with OpenACC if performance is so bad?

Analyzing OpenACC Run Time



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• The NVIDIA HPC SDK compiler provides automatic instrumentation when NV_ACC_TIME=1 at runtime

[alp514.hawk-b624](1002): NV_ACC_TIME=1 ./saxpyf_acc-nodata SAXPY Time: 0.778822

Accelerator Kernel Timina data /home/alp514/Workshop/2021HPC/parproa/solution/saxpy/saxpy acc-nodata.f90 saxov NVIDIA devicenum-0 time(us): 639.595 11: compute region reached 1 time 11: kernel launched 1 time arid: [65535] block: [128] device time(us): total=7.824 max=7.824 min=7.824 ava=7.824 elapsed time(us): total=7.872 max=7.872 min=7.872 ava=7.872 11: data region reached 2 times 16: data copyout transfers: 96 device time(us): total=243.594 max=2.558 min=1.751 ava=2.537 20: compute region reached 1 time 20: kernel launched 1 time arid: [65535] block: [128] device time(us): total=9.556 max=9.556 min=9.556 ava=9.556 elapsed time(us): total=9.605 max=9.605 min=9.605 ava=9.605 20: data region reached 2 times 20: data copvin transfers: 96 device time(us): total=256.886 max=2.709 min=1.841 ava=2.675 24: data copyout transfers: 48 device time(us): total=121.735 max=2.556 min=1.752 ava=2.536 [alp514.hawk-b624](1003): NV_ACC_TIME=1 ./saxpyc_acc-nodata SAXPY Time: 3.984324

Accelerator Kernel Timing data /home/alp514/Workshop/2021HPC/parproa/solution/saxpv/saxpv acc-nodata.c main NVIDIA devicenum-0 time(us): 1.277.749 16: compute region reached 1 time 16: kernel launched 1 time arid: [65535] block: [128] device time(us): total=14.513 max=14.513 min=14.513 ava=14.513 elapsed time(us): total=14.561 max=14.561 min=14.561 ava=14.561 16: data region reached 2 times 19: data copyout transfers: 192 device time(us): total=487.132 max=2.559 min=946 ava=2.537 22: compute region reached 1 time 22: kernel launched 1 time arid: [65535] block: [128] device time(us): total=19,484 max=19,484 min=19,484 ava=19,484 elapsed time(us): total=19.518 max=19.518 min=19.518 ava=19.518 22: data region reached 2 times 22: data copvin transfers: 192 device time(us): total=513.260 max=2.713 min=993 ava=2.673 24: data copyout transfers: 96 device time(us): total=243.360 max=2.556 min=944 ava=2.535

Fortran: ~ 17 ms for actual calculation and ~ 0.6 s for data transfer
C: ~ 35 ms for actual calculation and ~ 1.2 s for data transfer

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Processing Flow





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Processing Flow





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Processing Flow





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Defining data regions



• The data construct defines a region of code in which GPU arrays remain on the GPU and are shared among all kernels in that region



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- copy(list) Allocates memory on GPU and copies data from host to GPU when entering region and copies data to the host when exiting region.
- copyin(list) Allocates memory on GPU and copies data from host to GPU when entering region.
- copuout(list) Allocates memory on GPU and copies data to the host when exiting region.
 - create(list) Allocates memory on GPU but does not copy.
- present(list) Data is already present on GPU from another containing data region.
- Other clauses: present_or_copy[inlout], present_or_create, deviceptr.

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Array Shaping



• Compiler sometime cannot determine size of arrays

• Must specify explicitly using the data clauses and array "shape"

C #pragma acc data copyin(a[0:size]), copyout(b[s/4:3*s/4])
Fortran !\$acc data copyin(a(1:size)), copyout(b(s/4:3*s/4))

• Note: data clauses can be used on data, parallel or kernels



- Modify the SAXPY code to add a structured data region at the appropriate spot
- How does the compiler output the change?
- Is the code faster now?
- By how much and how does it compare with the serial and openmp code?
- Reprofile the code using NV_ACC_TIME

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Fortran Timings

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OpenMP (48 threads)	Xeon Gold 5220R	0.025263	19.97
OpenACC	Tesla T4	0.009601	52.55
CUDA	Tesla T4	0.007846	64.30

C Timings

Algorithm	Device	Time (s)	Speedup
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OpenMP (24 threads)	Xeon Gold 5220R	0.048442	10.57
OpenMP (48 threads)	Xeon Gold 5220R	0.026348	19.44
OpenACC	Tesla T4	0.019029	26.91

Update Construct



- Used to update existing data after it has changed in its corresponding copy (e.g. upate device copy after host copy changes).
- Move data from GPU to host, or host to GPU.
- Data movement can be conditional and asynchronous.
- Fortran

!\$acc update [clause ···]

• C

#pragma acc update [clause ···]

- Clause
 - host(list)
 - o device(list)
 - if(expression)
 - o async(expression)

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OpenACC private Clause



```
#pragma acc parallel loop
for(int i=0;i<M;i++) {
   for(int jj=0;jj<10;jj++)
      tmp[jj]=jj;
   int sum=0;
   for(int jj=0;jj<N;jj++)
      sum+=tmp[jj];
   A[i]=sum;
}</pre>
```

```
#pragma acc parallel loop
    private(tmp[0:9])
    for(int i=0;i<M;i++) {
        for(int jj=0;jj<10;jj++)
            tmp[jj]=jj;
        int sum=0;
        for(int jj=0;jj<N;jj++)
            sum+=tmp[jj];
        A[i]=sum;
    }
</pre>
```

• Compiler cannot parallelize because tmp is shared across threads

• Also useful for live-out scalars

```
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```

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OpenACC reduction Clause



• Reduction clause is allowed on *parallel* and *loop* constructs

Fortran

!\$acc parallel reduction(operation: var)
 structured block with reduction on var
!\$acc end parallel

C

#pragma acc kernels reduction(operation: var) {
 structured block with reduction on var
}

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Exercise



- Parallelize the Pi Calculation code using the Reduction clause
- Compare timing with serial and openmp code

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- Currently we have only exposed parallelism on the outer loop
- We know that both loops can be parallelized
- Lets look at methods for parallelizing multiple loops



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collapse(n): Applies the associated directive to the following ntightly nested loops.



Loops must be tightly nested

Exercise



- Parallelize the Matmul code using collapse clause
- Compare timings and GFLOPS with serial and openmp code

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Further Speedups



- OpenACC gives us more detailed control over parallelization
 - Via gang, worker and vector clauses
- By understanding more about specific GPU on which you're running, using these clauses may allow better performance.
- By understanding bottlenecks in the code via profiling, we can reorganize the code for even better performance.

General Principles: Finding Parallelism in Code



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- (Nested) for/do loops are best for parallelization
- Large loop counts are best
- Iterations of loops must be independent of each other
 - To help compiler: restrict keyword (C), independent clause
 - Use subscripted arrays, rather than pointer-indexed arrays
- Data regions should avoid wasted bandwidth
 - Can use directive to explicitly control sizes
- Various annoying things can interfere with accelerated regions.
 - Function calls within accelerated region must be inlineable.
 - No IO

OpenACC: Is it worth it?



- High-level. No involvement of OpenCL, CUDA, etc
- Single source. No forking off a separate GPU code. Compile the same program for accelerators or serial, non-GPU programmers can play along.
- Efficient. Experience shows very favorable comparison to low-level implementations of same algorithms.
- Performance portable. Supports GPU accelreators and co-processors from multiple vendors, current and future versions.
- Incremental. Developers can port and tune parts of their application as resources and profiling dictates. No wholesale rewrite required. Which can be quick.

Further Reading and References



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- OpenACC Programming and Best Practices Guide
- OpenACC 2.7 API Reference Card
- Parallel programming with OpenACC Rob Farber (Libraries Link)
- OpenACC for Programmers: Concepts and Strategies Guido Juckeland and Sunita Chandrasekaran (Libraries Link)

Lecture derived from slides and presentations by

- Michael Wolfe, PGI
- Jeff Larkin, NVIDIA
- John Urbanic, PSC

Search for OpenACC presentations at the GPU Technology Conference Website for further study http://www.gputechconf.com/gtcnew/on-demand-gtc.php

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